

Exploiting Cardinality Encodings in Parallel Maximum Satisfiability

Ruben Martins Vasco Manquinho Inês Lynce

IST/INESC-ID, Technical University of Lisbon, Portugal

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What is Maximum Satisfiability?

CNF Formula:

$$\begin{array}{cccc} x_6 \vee x_2 & \neg x_6 \vee x_2 & \neg x_2 \vee x_1 & \neg x_1 \\ \neg x_6 \vee x_8 & x_6 \vee \neg x_8 & x_2 \vee x_4 & \neg x_4 \vee x_5 \\ x_7 \vee x_5 & \neg x_7 \vee x_5 & \neg x_5 \vee x_3 & \neg x_3 \end{array}$$

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$\neg x_6 \vee x_8$	$x_6 \vee \neg x_8$	$x_2 \vee x_4$	$\neg x_4 \vee x_5$
$x_7 \vee x_5$	$\neg x_7 \vee x_5$	$\neg x_5 \vee x_3$	$\neg x_3$

- Formula is unsatisfiable

What is Maximum Satisfiability?

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$x_7 \vee x_5$	$\neg x_7 \vee x_5$	$\neg x_5 \vee x_3$	$\neg x_3$

- Formula is unsatisfiable
- Maximum Satisfiability (MaxSAT):
 - Find an assignment that maximizes (minimizes) number of satisfied (unsatisfied) clauses.

Motivation

- MaxSAT has several applications:
 - Software package upgradability;
 - Bug localization in C code;
 - Design debugging;
 - ...
- Improving MaxSAT solvers will have a practical impact;
- Multicore processors are now predominant;
- Parallel solving is known to boost the performance of the solver;

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- Improving MaxSAT solvers will have a practical impact;
- Multicore processors are now predominant;
- Parallel solving is known to boost the performance of the solver;
- **Goal:** develop a new parallel MaxSAT solver;
- **How:** using a portfolio of algorithms.

Algorithms for solving MaxSAT

Unsatisfiability-based algorithms:

$$x_6 \vee x_2 \quad \neg x_6 \vee x_2 \quad \neg x_2 \vee x_1 \quad \neg x_1$$

$$\neg x_6 \vee x_8 \quad x_6 \vee \neg x_8 \quad x_2 \vee x_4 \quad \neg x_4 \vee x_5$$

$$x_7 \vee x_5 \quad \neg x_7 \vee x_5 \quad \neg x_5 \vee x_3 \quad \neg x_3$$

Algorithms for solving MaxSAT

Unsatisfiability-based algorithms:

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$$\neg x_6 \vee x_8 \quad x_6 \vee \neg x_8$$

$$x_7 \vee x_5 \quad \neg x_7 \vee x_5$$

$$\neg x_2 \vee x_1 \quad \neg x_1$$

$$x_2 \vee x_4 \quad \neg x_4 \vee x_5$$

$$\neg x_5 \vee x_3 \quad \neg x_3$$

Formula is UNSAT; Get unsatisfiable sub-formula;
Lower bound value: 1

Algorithms for solving MaxSAT

Unsatisfiability-based algorithms:

$$x_6 \vee x_2$$

$$\neg x_6 \vee x_2$$

$$\neg x_2 \vee x_1 \vee r_1$$

$$\neg x_1 \vee r_2$$

$$\neg x_6 \vee x_8$$

$$x_6 \vee \neg x_8$$

$$x_2 \vee x_4 \vee r_3$$

$$\neg x_4 \vee x_5 \vee r_4$$

$$x_7 \vee x_5$$

$$\neg x_7 \vee x_5$$

$$\neg x_5 \vee x_3 \vee r_5$$

$$\neg x_3 \vee r_6$$

$$\sum_{i=1}^6 r_i \leq 1$$

Add relaxation variables and at-most-one constraint

Algorithms for solving MaxSAT

Unsatisfiability-based algorithms:

$$x_6 \vee x_2$$

$$\neg x_6 \vee x_2$$

$$\neg x_2 \vee x_1 \vee r_1$$

$$\neg x_1 \vee r_2$$

$$\neg x_6 \vee x_8$$

$$x_6 \vee \neg x_8$$

$$x_2 \vee x_4 \vee r_3$$

$$\neg x_4 \vee x_5 \vee r_4$$

$$x_7 \vee x_5$$

$$\neg x_7 \vee x_5$$

$$\neg x_5 \vee x_3 \vee r_5$$

$$\neg x_3 \vee r_6$$

$$\sum_{i=1}^6 r_i \leq 1$$

Formula is UNSAT; Get unsatisfiable sub-formula;

Lower bound value: 2

Algorithms for solving MaxSAT

Unsatisfiability-based algorithms:

$$\begin{array}{cccc} x_6 \vee x_2 \vee r_7 & \neg x_6 \vee x_2 \vee r_8 & \neg x_2 \vee x_1 \vee r_1 \vee r_9 & \neg x_1 \vee r_2 \vee r_{10} \\ \neg x_6 \vee x_8 & x_6 \vee \neg x_8 & x_2 \vee x_4 \vee r_3 & \neg x_4 \vee x_5 \vee r_4 \\ x_7 \vee x_5 \vee r_{11} & \neg x_7 \vee x_5 \vee r_{12} & \neg x_5 \vee x_3 \vee r_5 \vee r_{13} & \neg x_3 \vee r_6 \vee r_{14} \\ \sum_{i=1}^6 r_i \leq 1 & \sum_{i=7}^{14} r_i \leq 1 & & \end{array}$$

Add relaxation variables and at-most-one constraint

Algorithms for solving MaxSAT

Unsatisfiability-based algorithms:

$$\begin{array}{cccc} x_6 \vee x_2 \vee r_7 & \neg x_6 \vee x_2 \vee r_8 & \neg x_2 \vee x_1 \vee r_1 \vee r_9 & \neg x_1 \vee r_2 \vee r_{10} \\ \neg x_6 \vee x_8 & x_6 \vee \neg x_8 & x_2 \vee x_4 \vee r_3 & \neg x_4 \vee x_5 \vee r_4 \\ x_7 \vee x_5 \vee r_{11} & \neg x_7 \vee x_5 \vee r_{12} & \neg x_5 \vee x_3 \vee r_5 \vee r_{13} & \neg x_3 \vee r_6 \vee r_{14} \\ \sum_{i=1}^6 r_i \leq 1 & \sum_{i=7}^{14} r_i \leq 1 & & \end{array}$$

Formula is now SAT

The minimum number of unsatisfiable clauses is 2

Algorithms for solving MaxSAT

Linear search algorithms on the number of unsatisfiable clauses:

$x_6 \vee x_2$	$\neg x_6 \vee x_2$	$\neg x_2 \vee x_1$	$\neg x_1$
$\neg x_6 \vee x_8$	$x_6 \vee \neg x_8$	$x_2 \vee x_4$	$\neg x_4 \vee x_5$
$x_7 \vee x_5$	$\neg x_7 \vee x_5$	$\neg x_5 \vee x_3$	$\neg x_3$

Algorithms for solving MaxSAT

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Formula is UNSAT

Algorithms for solving MaxSAT

Linear search algorithms on the number of unsatisfiable clauses:

$$x_6 \vee x_2 \vee r_1$$

$$\neg x_6 \vee x_2 \vee r_2$$

$$\neg x_2 \vee x_1 \vee r_3$$

$$\neg x_1 \vee r_4$$

$$\neg x_6 \vee x_8 \vee r_5$$

$$x_6 \vee \neg x_8 \vee r_6$$

$$x_2 \vee x_4 \vee r_7$$

$$\neg x_4 \vee x_5 \vee r_8$$

$$x_7 \vee x_5 \vee r_9$$

$$\neg x_7 \vee x_5 \vee r_{10}$$

$$\neg x_5 \vee x_3 \vee r_{11}$$

$$\neg x_3 \vee r_{12}$$

$$\sum_{i=1}^{12} r_i \leq 12$$

Add relaxation variables and at-most-k constraint

Algorithms for solving MaxSAT

Linear search algorithms on the number of unsatisfiable clauses:

$$x_6 \vee x_2 \vee r_1$$

$$\neg x_6 \vee x_2 \vee r_2$$

$$\neg x_2 \vee x_1 \vee r_3$$

$$\neg x_1 \vee r_4$$

$$\neg x_6 \vee x_8 \vee r_5$$

$$x_6 \vee \neg x_8 \vee r_6$$

$$x_2 \vee x_4 \vee r_7$$

$$\neg x_4 \vee x_5 \vee r_8$$

$$x_7 \vee x_5 \vee r_9$$

$$\neg x_7 \vee x_5 \vee r_{10}$$

$$\neg x_5 \vee x_3 \vee r_{11}$$

$$\neg x_3 \vee r_{12}$$

$$\sum_{i=1}^{12} r_i \leq 12$$

Formula is SAT; solution found has $\sum_{i=1}^{12} r_i = 9$;
Upper bound value: 9

Algorithms for solving MaxSAT

Linear search algorithms on the number of unsatisfiable clauses:

$$x_6 \vee x_2 \vee r_1$$

$$\neg x_6 \vee x_2 \vee r_2$$

$$\neg x_2 \vee x_1 \vee r_3$$

$$\neg x_1 \vee r_4$$

$$\neg x_6 \vee x_8 \vee r_5$$

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$$\neg x_7 \vee x_5 \vee r_{10}$$

$$\neg x_5 \vee x_3 \vee r_{11}$$

$$\neg x_3 \vee r_{12}$$

$$\sum_{i=1}^{12} r_i \leq 8$$

Update at-most-k constraint

Algorithms for solving MaxSAT

Linear search algorithms on the number of unsatisfiable clauses:

$$x_6 \vee x_2 \vee r_1$$

$$\neg x_6 \vee x_2 \vee r_2$$

$$\neg x_2 \vee x_1 \vee r_3$$

$$\neg x_1 \vee r_4$$

$$\neg x_6 \vee x_8 \vee r_5$$

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$$x_7 \vee x_5 \vee r_9$$

$$\neg x_7 \vee x_5 \vee r_{10}$$

$$\neg x_5 \vee x_3 \vee r_{11}$$

$$\neg x_3 \vee r_{12}$$

$$\sum_{i=1}^{12} r_i \leq 8$$

Formula is SAT; solution found has $\sum_{i=1}^{12} r_i = 2$;
Upper bound value: 2

Algorithms for solving MaxSAT

Linear search algorithms on the number of unsatisfiable clauses:

$$x_6 \vee x_2 \vee r_1$$

$$\neg x_6 \vee x_2 \vee r_2$$

$$\neg x_2 \vee x_1 \vee r_3$$

$$\neg x_1 \vee r_4$$

$$\neg x_6 \vee x_8 \vee r_5$$

$$x_6 \vee \neg x_8 \vee r_6$$

$$x_2 \vee x_4 \vee r_7$$

$$\neg x_4 \vee x_5 \vee r_8$$

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$$\sum_{i=1}^{12} r_i \leq 1$$

Update at-most-k constraint

Algorithms for solving MaxSAT

Linear search algorithms on the number of unsatisfiable clauses:

$$x_6 \vee x_2 \vee r_1$$

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$$\neg x_3 \vee r_{12}$$

$$\sum_{i=1}^{12} r_i \leq 1$$

Formula is UNSAT;

The minimum number of unsatisfiable clauses is 2

Cardinality constraints in MaxSAT solving

- MaxSAT algorithms use cardinality constraints:
 - at-most-one;
 - at-most-k.
- Cardinality constraints:
 - $\sum_{i=1}^n l_i \leq k$, where $l_i \in \{x_i, \neg x_i\}$
- Handling cardinality constraints:
 - Encode cardinality constraints into clauses;
 - Use a native representation of cardinality constraints.

Encodings for cardinality constraints

- Cardinality constraints:
 - At-most-one
 - At-most-k

Encodings for cardinality constraints

- Cardinality constraints:
 - At-most-one
 - At-most-k
- Naive (pairwise) encoding for at-most-one constraints:
 - Cardinality constraint: $x_1 + x_2 + x_3 + x_4 \leq 1$
 - Clauses:

$$\left. \begin{array}{l} (x_1 \Rightarrow \neg x_2) \\ (x_1 \Rightarrow \neg x_3) \\ (x_1 \Rightarrow \neg x_4) \\ \dots \end{array} \right\} \begin{array}{l} \neg x_1 \vee \neg x_2 \\ \neg x_1 \vee \neg x_3 \\ \neg x_1 \vee \neg x_4 \\ \dots \end{array}$$

- Complexity: $\mathcal{O}(n^2)$ clauses

Encodings for cardinality constraints

- Cardinality constraints:
 - At-most-one
 - At-most-k
- Naive encoding for at-most-k constraints:
 - Cardinality constraint: $x_1 + x_2 + x_3 + x_4 \leq 2$
 - Clauses:

$$\left. \begin{array}{l} (x_1 \wedge x_2 \Rightarrow \neg x_3) \\ (x_1 \wedge x_2 \Rightarrow \neg x_4) \\ (x_2 \wedge x_3 \Rightarrow \neg x_4) \\ \dots \end{array} \right\} \begin{array}{l} (\neg x_1 \vee \neg x_2 \vee \neg x_3) \\ (\neg x_1 \vee \neg x_2 \vee \neg x_4) \\ (\neg x_2 \vee \neg x_3 \vee \neg x_4) \\ \dots \end{array}$$

- Complexity: $\mathcal{O}(n^k)$ clauses

Encodings for cardinality constraints

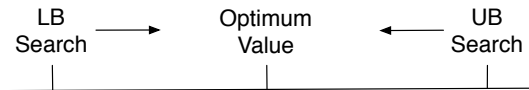
Encoding	Clauses	Variables	Type
Pairwise	$\mathcal{O}(n^2)$	0	at-most-one
Ladder	$\mathcal{O}(n)$	$\mathcal{O}(n)$	at-most-one
Bitwise	$\mathcal{O}(n \log_2 n)$	$\mathcal{O}(\log_2 n)$	at-most-one
Commander	$\mathcal{O}(n)$	$\mathcal{O}(n)$	at-most-one
Product	$\mathcal{O}(n)$	$\mathcal{O}(n)$	at-most-one
Sequential	$\mathcal{O}(nk)$	$\mathcal{O}(nk)$	at-most-k
Totalizer	$\mathcal{O}(nk)$	$\mathcal{O}(n \log_2 n)$	at-most-k
Sorters	$\mathcal{O}(n \log_2^2 n)$	$\mathcal{O}(n \log_2^2 n)$	at-most-k

Encodings for cardinality constraints

- Cardinality constraints:
 - At-most-one;
 - At-most-k.
- Many encodings for cardinality constraints are available;
- Different encodings can perform better over different problems;
- Exploit the diversification of cardinality encodings.

Parallel MaxSAT solver (2 threads)

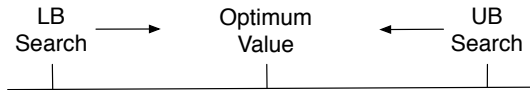
- Search in the lower and upper bound values of the optimal solution:



- The optimum value is found when:
 - LB or UB search terminates with a solution;
 - or when LB value = UB value.

Parallel MaxSAT solver (2 threads)

- Search in the lower and upper bound values of the optimal solution:

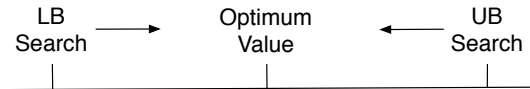


- The optimum value is found when:
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 - or when LB value = UB value.

T_0													T_1
LB													UB
0	1	2	3	4	5	6	7	8	9	10	11	12	

Parallel MaxSAT solver (2 threads)

- Search in the lower and upper bound values of the optimal solution:



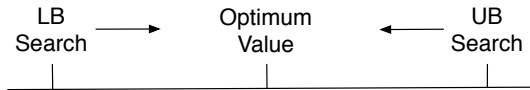
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T_0 returns UNSAT; update lower bound value

Parallel MaxSAT solver (2 threads)

- Search in the lower and upper bound values of the optimal solution:



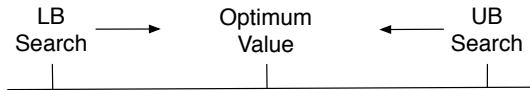
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	T_0								T_1			
	LB								UB			
0	1	2	3	4	5	6	7	8	9	10	11	12

T_1 returns SAT; update upper bound value

Parallel MaxSAT solver (2 threads)

- Search in the lower and upper bound values of the optimal solution:



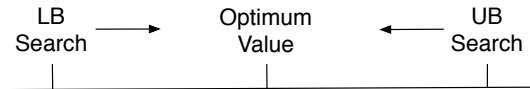
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		T_0							T_1			
		LB							UB			
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Parallel MaxSAT solver (2 threads)

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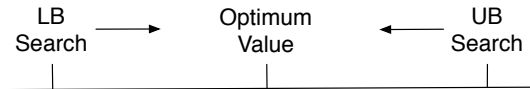
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		$T_0; T_1$										
		LB; UB										
0	1	2	3	4	5	6	7	8	9	10	11	12

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Parallel MaxSAT solver (2 threads)

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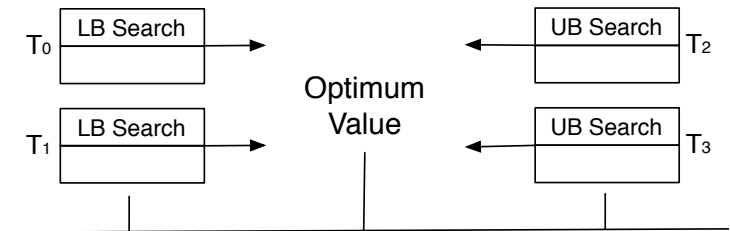
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		$T_0; T_1$											
		LB; UB											
0	1	2	3	4	5	6	7	8	9	10	11	12	

LB value = UB value, optimal value has been found

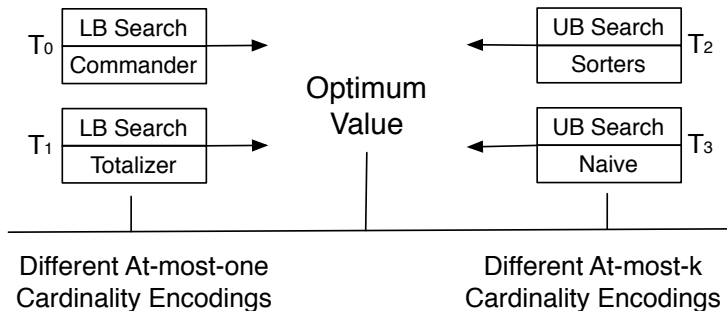
Parallel MaxSAT solver (n threads)

- Search in the lower and upper bound values of the optimal solution:



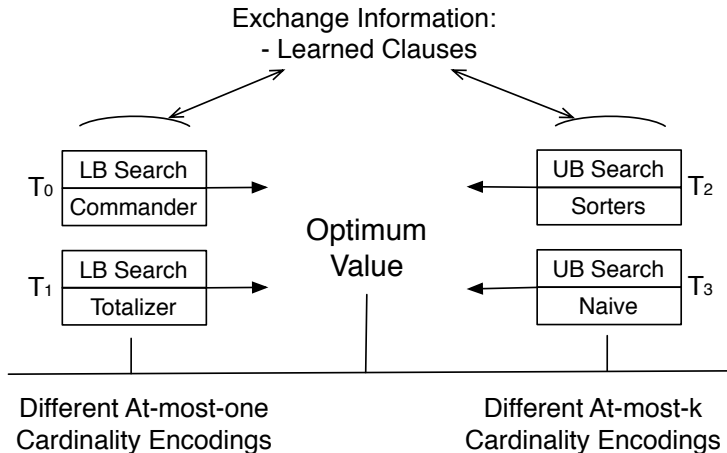
Parallel MaxSAT solver (n threads)

- Search in the lower and upper bound values of the optimal solution:



Parallel MaxSAT solver (n threads)

- Search in the lower and upper bound values of the optimal solution:



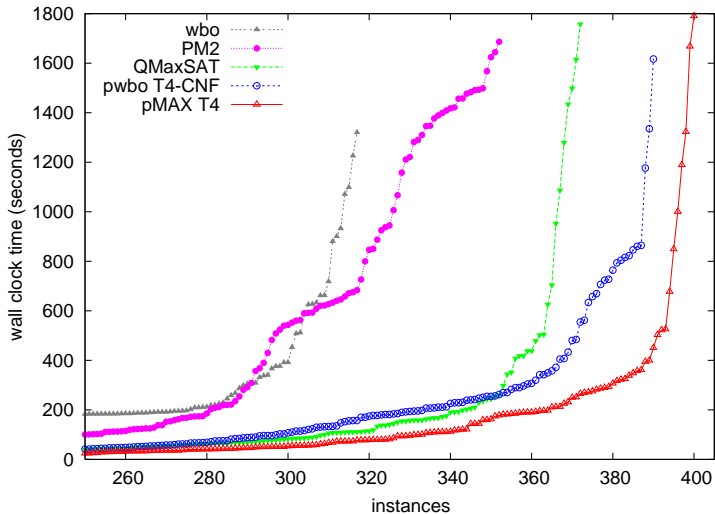
Experimental Results

- Sequential Solvers:
 - QMaxSAT:
 - Upper bound search solver;
 - Uses cardinality encodings.
 - PM2:
 - Lower bound search solver;
 - Uses cardinality encodings.
 - wbo:
 - Upper bound search during 10% of the time limit;
 - Lower bound search on remaining time;
 - Does not use cardinality encodings.

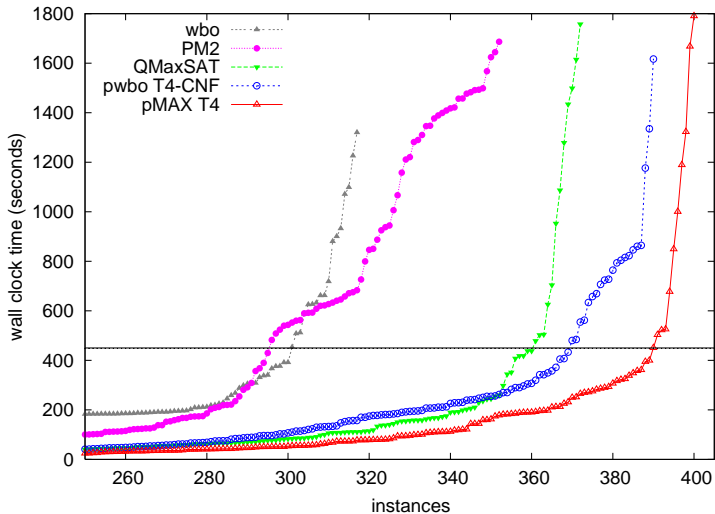
Experimental Results

- Parallel Solvers:
 - pwbo-CNF:
 - Searches on the lower and upper bound values of the optimal solution;
 - Splits the search on different upper bound values;
 - Only one thread uses cardinality encodings.
 - pMAX:
 - The solver proposed in this talk;
 - Searches on the lower and upper bound values of the optimal solution;
 - Uses a different cardinality encoding in each thread.

Experimental Results

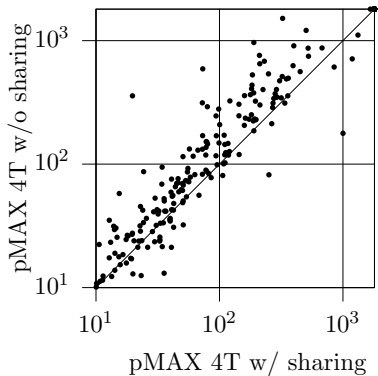


Experimental Results



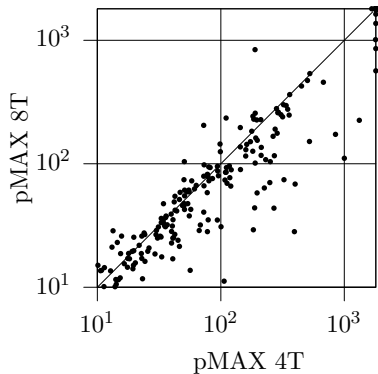
Experimental Results

- Impact of sharing learned clauses (seconds):



Experimental Results

- Scalability of pMAX (seconds):



Experimental Results

- Speedup on instances solved by all our solvers:

Solver	Time (s)	Speedup
wbo	67,947.41	1.00
pwbo 4T-CNF	18,015.69	3.77
pMAX 4T	11,382.91	5.97
pMAX 8T	7,990.10	8.50

Conclusions

- Diversification of cardinality encodings can be used in parallel MaxSAT;
- pMAX outperforms state-of-the-art MaxSAT solvers:
 - Even when considering CPU time.
- Sharing learned clauses has a strong impact on the solving speed;
- pMAX shows scalability:
 - pMAX 8T is $1.4\times$ faster than pMAX 4T.