Towards Smarter MACE-style Model Finders

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Background and Motivation



Does a FOL formula have a finite model?

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- Finite models useful:
 - "debugging" of wrong theorems
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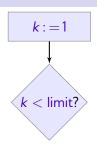
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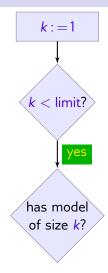
Advatage:

- If a finite model exists, it is fount in finite time.
- Complete for some theories (Bernays-Schönfinkel, a.k.a. EPR)

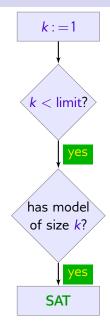




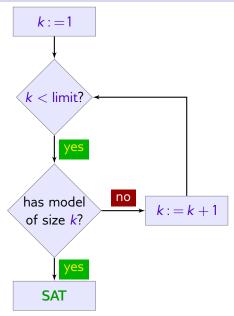




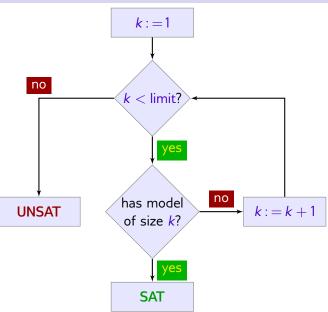












Avoiding Space Explosion



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Encoding directly to SAT is exponential, eventually blows up

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Remedy:

Encode into SAT lazily by

Count-Eexample Abstraction Guided Refinement (CEGAR)



$$(\exists \vec{p} \, \vec{f} \,)(\forall \vec{x} \,) \, \phi$$

 \vec{p} predicates, \vec{f} functions, \vec{x} FOL variables

$$\alpha := true$$



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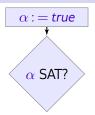
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- **7 GOTO** 2

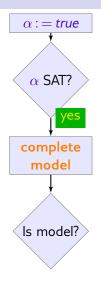




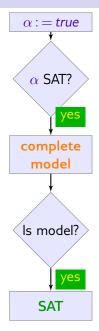




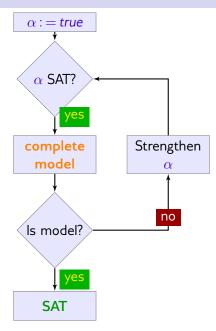




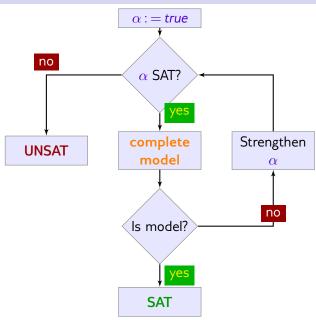












Completing Models



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$$p(0), \neg q(0) \models p(0) \lor q(0)$$

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We need to complete into **interpretation** of original **Examples** for $(\forall x)(p(x) \lor \neg q(y))$

$$p \triangleq \{0\}, q \triangleq \{\}$$

$$p \triangleq \{0\}, q \triangleq \{1\}$$

$$p \triangleq 2^{1..k}, q \triangleq \{\}$$

Completing Models Contd.



Natural approach: set undefined to false/true **Examples**

$$\{p(0), \neg p(1), \neg p(2)\} \dots p \triangleq 2^{1..k} \setminus \{1, 2\}$$

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Completing Models Contd.



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Remedy:

Learn the completion with Machine Learning techniques.





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$$t \triangleq 1$$

$$p(0,\ldots,0) \triangleq \mathsf{False}$$

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- **5** Model of α :

$$t \triangleq 1$$

 $p(0, ..., 0) \triangleq \text{False}$
 $p(1, ..., 0) \triangleq \text{True}$

6 Learn:

$$t \triangleq 1$$

$$p(x_1,\ldots,x_n)\triangleq(x_1=1)$$

Technical Remarks (QFM)



■ Learning by decision trees



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- Function and predicates eliminated by Ackermann reduction



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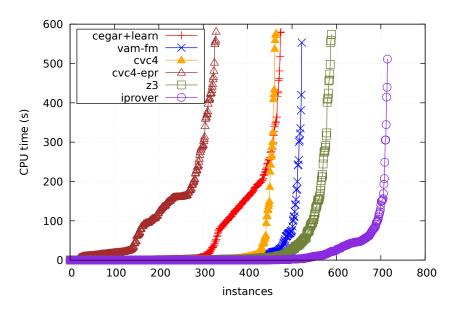
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- Incremental SAT (minisat)
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- Symmetry breaking, e.g. $c_1 \triangleq 0$

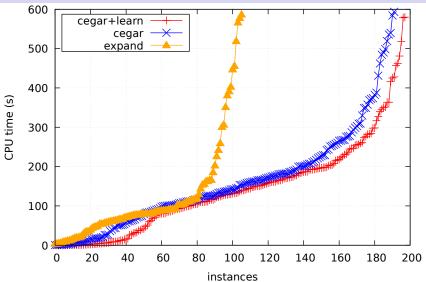
Results EPR





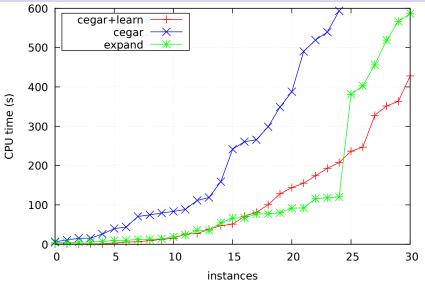
Results EPR: QFM





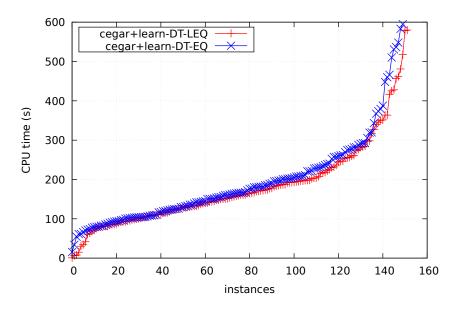
Results EPR QFM (SAT)





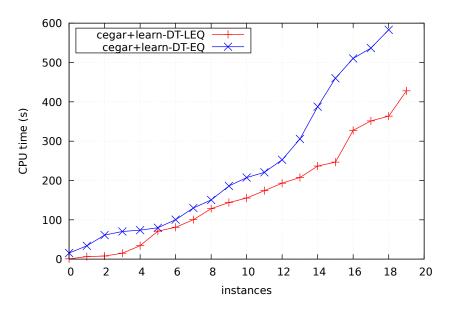
Results EPR: Learning Method





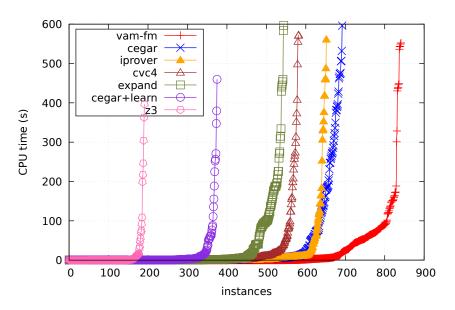
Results EPR: Learning Method (SAT)





Results SAT NON-EPR







■ CEGAR for lazy SAT-based model finite model finding



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- Observing a formula while solving, learn from that
- Better learning methods?
- Learning in the presence of theories?
- Infinite domains?



http://sat2019.tecnico.ulisboa.pt



